

# Graph Theory Preliminary Exam

26 June 1999

**Instructions:** Do all four problems in **Part A** and do exactly four of the five problems in **Part B**. Indicate clearly which problem in Part B you have omitted. Each problem in Part A is valued at 10 points, while each problem in Part B is valued at 15 points. Hand in **eight** problems only. Begin each problem on a new sheet of paper and only write on one side of the paper. You have six hours to complete the exam. When you are ready to hand in your exam, assemble your solutions in numerical order, write your name on the front page, and initial all other pages.

**Part A: Do all four problems.** (10 points for each problem)

- A1 Prove that every nontrivial connected graph contains at least two vertices that are not cut-vertices.
- A2 Let  $G = K_{3,3} - e$ , where  $K_{3,3}$  has partite sets  $V_1 = \{v_1, v_3, v_5\}$  and  $V_2 = \{v_2, v_4, v_6\}$  and  $e = v_5v_6$ .
- Describe a 2-cell embedding of  $G$  on a surface  $S_k$ , where  $k \neq \text{gen}(G)$ , by giving an appropriate 6-tuple of cyclic permutations. What is  $k$  in your example? Include all details.
  - On what surfaces is  $G$  2-cell embeddable? Verify your answer.
- A3 The *subdivision graph*  $S(G)$  of a graph  $G$  is that graph obtained from  $G$  by replacing each edge  $uv$  of  $G$  by a new vertex  $w$  and the two edges  $uw$  and  $vw$ . Prove that if  $S(G)$  is hamiltonian, then  $G$  is eulerian.
- A4 Recall that  $\omega(G)$  denotes the clique number of  $G$  (the maximum number of vertices in a complete subgraph of  $G$ ). Certainly,  $\chi(G) \geq \omega(G)$  for every graph  $G$ . Define a graph  $G$  to be *perfect* if  $\chi(F) = \omega(F)$  for every induced subgraph  $F$  of  $G$  (including  $F = G$ ).
- Explain why every complete graph and every bipartite graph is perfect.
  - Prove that a critically  $k$ -chromatic graph  $G$  is perfect if and only if  $G = K_k$ .

**Part B: Do exactly four of the five problems.** (15 points for each problem)

B1 Let  $T_n$  denote the transitive tournament of order  $n$ . Hence  $T_n$  can be defined as that digraph with vertex set  $\{v_1, v_2, \dots, v_n\}$  such that  $(v_i, v_j)$  is an arc of  $T_n$  if and only if  $i < j$ . Also, recall that an oriented graph  $G$  is obtained by assigning a direction to each edge of some graph  $G$ . For oriented graphs  $D_1$  and  $D_2$ , define the *tournament Ramsey number*  $TR(D_1, D_2)$  as the smallest positive integer  $p$  for which any 2-coloring of the arcs of  $T_p$  results in either a red  $D_1$  or a blue  $D_2$ .

- (a) Determine  $TR(T_2, T_m)$ .
- (b) Find necessary and sufficient conditions on  $D_1$  and  $D_2$  in order to guarantee the existence of  $TR(D_1, D_2)$ .
- (c) Let  $\vec{P}_{n+1}$  denote the directed path of order  $n + 1$ . Prove that

$$TR(\vec{P}_{n+1}, \vec{P}_{m+1}) = nm + 1.$$

[Hint: Let  $r_i$  denote the maximum order of a red directed path starting at vertex  $v_i$  and let  $b_i$  denote the maximum order of a blue directed path starting at  $v_i$ .]

B2 Let the graph  $G$  be obtained from  $K_{2,6}$ , whose partite sets are  $V_1 = \{u_1, u_2\}$  and  $V_2 = \{v_1, v_2, \dots, v_6\}$ , by adding two new vertices  $w_1, w_2$  and the seven edges  $u_1u_2, w_1v_i$  ( $1 \leq i \leq 3$ ), and  $w_2v_j$  ( $4 \leq j \leq 6$ ). The graph  $G$  is shown below.

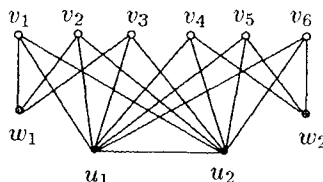


Figure 1: The graph  $G$

- (a) Use Tutte's characterization of graphs with 1-factors to show that  $G$  does not have a 1-factor.
- (b) Petersen's theorem states that if  $G$  is a bridgeless, cubic graph, then  $G$  has a 1-factor. Show that Petersen's theorem can be extended somewhat by proving that if  $G$  is a bridgeless graph with degree set  $\{3, 7\}$  such that  $G$  has exactly one vertex of degree 7, then  $G$  has a 1-factor.
- (c) Show that the result in (b) cannot be extended any further by giving an example of a bridgeless graph  $G$  with degree set  $\{3, 7\}$  such that  $G$  has two vertices of degree 7 and no 1-factor.